



Queens of the Hill

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Core Wars



Core Wars is a programming game in which two or more programs run in a simulated computer with the goal of terminating every other program and surviving as long as possible.

corewars.org

MARS: the Core Wars environment

Common memory space

Only instructions

↳ data is part of the instructions

Language: Redcode

...
MOV 0, 1
...
ADD #4, 3
MOV 2, @2
JMP -2
DAT #0, #0
...

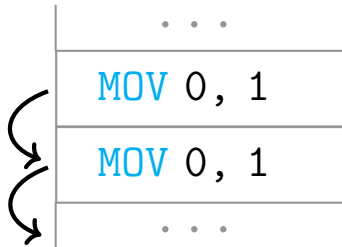
The Imp: the simplest warrior

MOV 0, 1

Relative memory addresses:

↳ 0 current, 1 next.

Program: copy the current instruction to the next address.



Result: The Imp copies itself all over memory.

The Imp does not win

The Imp is **good at survival**, but **bad at killing**.

↳ it kills no processes.

The **DAT** instruction kills the current process.

To kill a process, insert **DAT** and make it execute it.

The Dwarf

Bomb the memory with regularly spaced **DAT**.

- | | | |
|----|-------------------|---|
| 0: | ADD #4, 3 | Add 4 to instruction 3.
↳ its 2nd argument |
| 1: | MOV 2, @2 | Move instruction 1+2 to the
value of its @2nd argument. |
| 2: | JMP -2 | Jump back 2 instructions. |
| 3: | DAT #0, #0 | |

The Dwarf

Bomb the memory with regularly spaced **DAT**.

0: **ADD** #4, 3

Add 4 to instruction 3.

↳ its 2nd argument

1: **MOV** 2, @2

Move instruction 1+2 to the value of its @2nd argument.

2: **JMP** -2

Jump back 2 instructions.

3: **DAT** #0, #4

4: . . .

5: . . .

6: . . .

7: **DAT** #0, #4

8: . . .

} **DAT** at addresses not dangerous to the Dwarf.

Core Wars Tournaments

King of the hill mode:

- 10–30 warriors
- sequentially interleaving runs
- score = $f(\text{number of killed rivals})$

Highest score: current king of the hill

Lowest score:

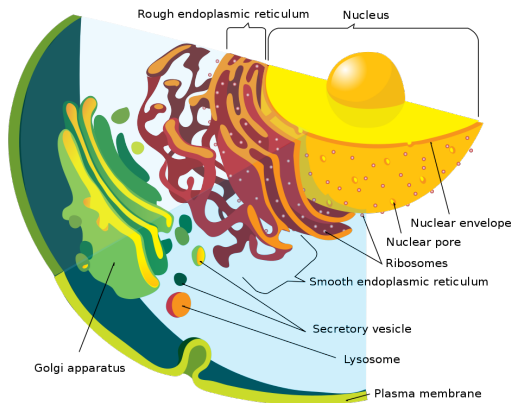
- 1 push off the hill
- 2 replace by the next contestant

So what?

P systems vs. Life

Inspired by the eukaryotic cell

Decentralized computing



$$a \rightarrow aa$$

$$a \rightarrow (a, \text{out})$$

a

1

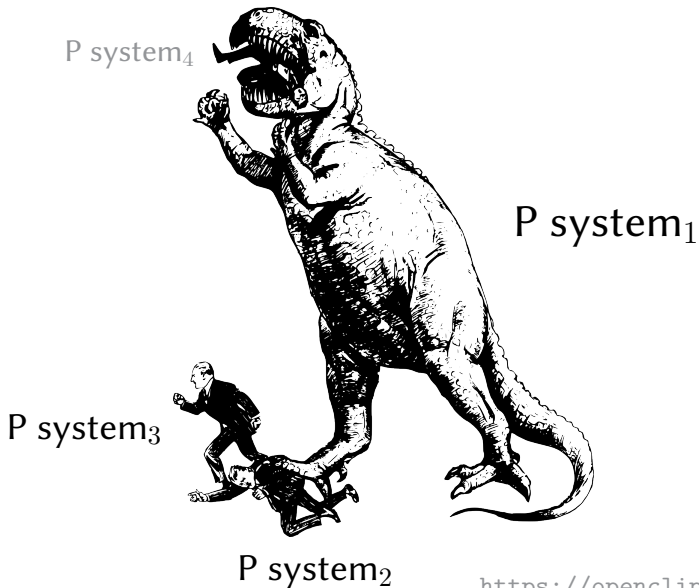
$$a \rightarrow b$$

$$b \rightarrow c (c, \text{in})$$

0

Use P systems as a **tool for thinking** about Life.

P systems vs. Evolution



<https://openclipart.org/>



Queens of the Hill

Run tournaments between P systems

|
Valkyries

Valkyries: expectations of the formalism

- Ease of **interaction** in a group of valkyries.
- Ease of **programming** individual valkyries.



Transition P systems with communication and anti-matter.

Valkyrie P systems

$$\Pi = (O, \mu, w_1, \dots, w_n, R_1, \dots, R_n)$$

- $O = \Sigma \cup \Delta_k$: the finite **alphabet** of objects
- $\Delta_k = \{\delta_t, \bar{\delta}_t \mid 1 \leq t \leq k\} \cup \{\delta\}$, $k \in \mathbb{N}$:
the **dissolution timers**
- μ : the hierarchical **membrane structure**
- w_i : the **initial multiset** in membrane i
- R_i : the finite **set of rules** in membrane i

Rule types

- 1 Full cooperation: $abc \rightarrow xyz$
- 2 Target indications: $Tar = \{in, here, out\}$:
 $ab \rightarrow (c, out)(d, in)$
- 3 Membrane dissolution: $ab \rightarrow \delta$
 - ▶ all objects and the inner membranes fall through to the parent membrane
- 4 Dissolution timers: $\delta_t \rightarrow \delta_{t-1}, \delta_1 \rightarrow \delta$
- 5 Anti-matter annihilation for Δ_k : $\delta_t \bar{\delta}_t \rightarrow \lambda$.
 - ▶ weak priority for annihilation
 - ▶ Δ_k **forbidden** in normal LHS.

Computations and **not** halting

Usual semantics:

- 1 Apply the rules in the **maximally parallel** manner.
- 2 Perform the **dissolutions**.

Halting configurations: no more applicable rules.

Don't care about halting: **continuous computation**

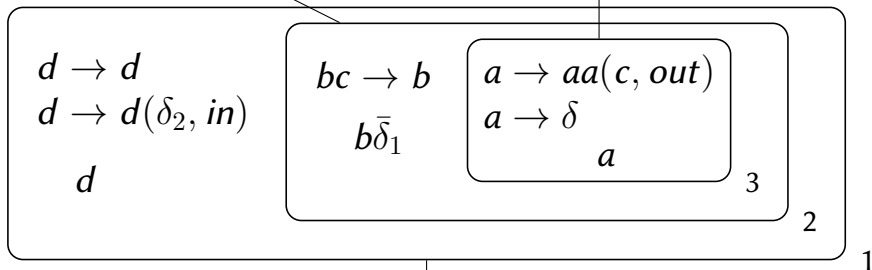
↳ like in VAS, Lindenmayer systems, and Core Wars

Example

annihilation rules not shown

1

- double some a and eject c
- dissolve the membrane
- b progressively erases c

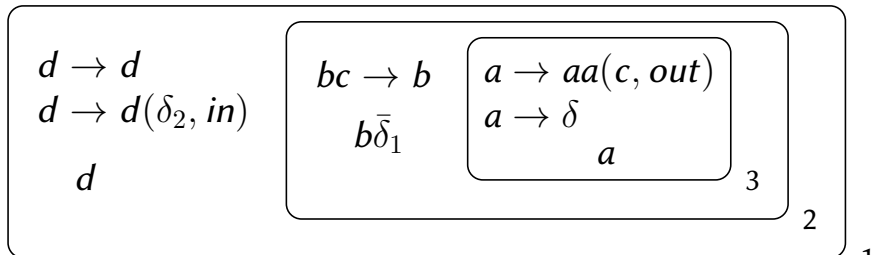


- maintain d
- maintain d and inject δ_2

Example

annihilation rules not shown

2

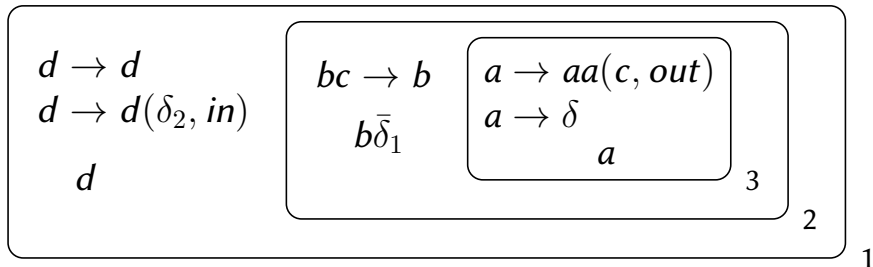


- First δ_2 in membrane 2 will become δ_1 and annihilate with $\bar{\delta}_1$.
- Second δ_2 will dissolve membrane 2.

Example

annihilation rules not shown

3



Possible evolutions:

		w_1	w_2	w_3
dissolved	\emptyset	d	$b\bar{\delta}_1 c^*$	a^{2^k}
	2	$dbc^* \{\delta_2, \delta_1, \lambda\}$		a^{2^k}
	3	d	$b\bar{\delta}_1 c^* a^*$	
	2, 3	$dbc^* a^* \{\delta_2, \delta_1, \lambda\}$		

Tournaments

$$Q = \left(\begin{array}{ccc} \boxed{\Pi_1} & \cdots & \boxed{\Pi_m} \\ (1, 1) & & (m, 1) \\ \{a \rightarrow (a, in) \mid a \in O\} \end{array} \right)_0$$

- **Common alphabet** $O = \Sigma \cup \Delta_k$ for all Π_i .
- **Membrane i in Π_j** \rightsquigarrow membrane (j, i) .

Tournament semantics

Same semantics for Q as for individual valkyries.

Dissolving the skin of a valkyrie is allowed \Rightarrow **killing**.

The skin **bounces all symbols back**.

\hookrightarrow including δ_k and $\bar{\delta}_k$

The symbols may end up in another valkyrie

\Rightarrow **communication by non-determinism**.

Tournament organization

- 1 Run all valkyries in *max* mode for N steps, resolving non-determinism **probabilistically**.
- 2 Repeat 1 M times.
- 3 Compute the **score** of the valkyrie Π_j based on how many of its membranes were dissolved.

$$\text{score}(\Pi_j) = \frac{1}{|\Pi_j|} \left(|\Pi_j| - \frac{1}{M} \sum_{i=1}^M \text{diss}_i(\Pi_j) \right)$$

Tournament scoring

$$\text{score}(\Pi_j) = \frac{1}{|\Pi_j|} \left(|\Pi_j| - \frac{1}{M} \sum_{i=1}^M \text{diss}_i(\Pi_j) \right)$$

- $\text{diss}_i(\Pi_j)$: number of membranes of Π_j that were dissolved during i -th run of ①
- $|\Pi_j|$: total number of membranes in Π_j

Computing the score: an example

$$\text{score}(\Pi_j) = \frac{1}{|\Pi_j|} \left(|\Pi_j| - \frac{1}{M} \sum_{i=1}^M \text{diss}_i(\Pi_j) \right)$$

Let $|\Pi_j| = 5$, $M = 3$, and suppose 2, 3, and 4 membranes were dissolved:

$$\text{diss}_1(\Pi_j) = 2 \quad \text{diss}_2(\Pi_j) = 3 \quad \text{diss}_3(\Pi_j) = 4$$

$$\text{score}(\Pi_j) = \frac{1}{5} \left(5 - \frac{2 + 3 + 4}{3} \right) = \frac{2}{5}$$

Tournament parameters

m	10–20	The number of entrants.
N	1000	The length of a computation.
M	50	The total number of computations.
k	5	The maximal value of the index t in δ_t .
$ \Sigma $	10	The number of working symbols.

Values derived from similar experiments with multi-agent systems \Rightarrow **Test and improve!**

Computational complexity

Quickly.

Valkyries are computationally complete

Simulate $(p, \text{ADD}(r), q)$ by $p \rightarrow qa_r$.

Simulate $(p, \text{SUB}(r), q, s)$:

Decrement

1. $p \rightarrow \bar{p}_1 \hat{p}_1$

2. $\bar{p}_1 a_r \rightarrow \bar{p}_2, \hat{p}_1 \rightarrow \hat{p}_2$

3. $\hat{p}_2 \bar{p}_2 \rightarrow q, \hat{p}_2 \bar{p}_1 \rightarrow \#$

Zero test

$p \rightarrow \tilde{p}_1 \dot{p}_1$

$\dot{p}_1 a_r \rightarrow \#, \tilde{p}_1 \rightarrow \tilde{p}_2$

$\tilde{p}_2 \dot{p}_1 \rightarrow s$

The language is **rich enough**.

Tournaments are **not** computations

The proof relies on **non-determinism** and **halting**.

No halting in tournaments.

The non-determinism is resolved **probabilistically**.

⇒ **Partial biased** coverage of the computation tree.

Efficiency > Expressive power

First valkyries

The Bomber

Bomb around with δ_t .

$$\{a \rightarrow a(\delta_t, \text{out}) \mid 1 \leq t \leq k\}$$

a

1

Number of valkyries \downarrow Efficiency \downarrow

- With 2 valkyries, δ_t may come back into the Bomber.

The Bar Bomber

Bomb around with δ_t , but also **stock up** $\bar{\delta}_t$.

$$\{a \rightarrow a \bar{\delta}_t (\delta_t, out) \mid 1 \leq t \leq k\}$$

a

1

Protects against other Bombers.

Overwhelmed when too many Bombers around.

The Anti-Bomber

Bomb around with δ_t , but also eject $\bar{\delta}_t$.

$$\{a \rightarrow a(\bar{\delta}_t, \text{out})(\delta_t, \text{out}) \mid 1 \leq t \leq k\}$$

a

1

Protects this valkyrie, but also the other valkyries.

The Delta Wall

If the number of entrants m is known:

↳ or the upper bound on m

$$a \rightarrow a\bar{\delta}_1^{r(m-1)}$$

$$a\bar{\delta}_1^{r(m-1)}$$

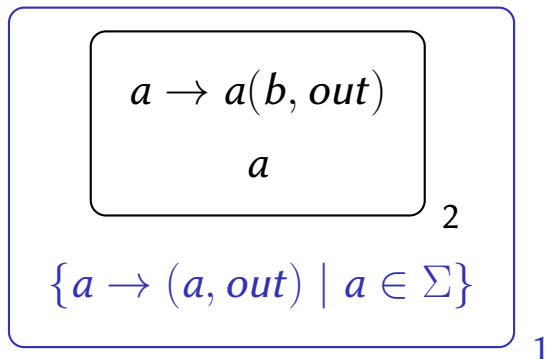
1

Stock up “enough” copies of $\bar{\delta}_1$.

↳ r defines “enough”

The 2-layer Onion

Wrap the core valkyrie in layers.



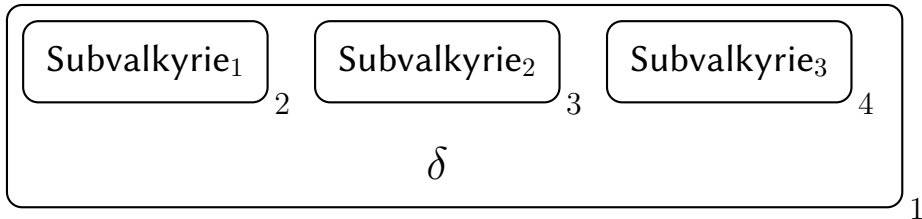
How to emit δ_t without dissolving membrane 1?

Destabilize other valkyries by pushing other symbols.

The Bombshell

Release **multiple valkyries** from a common skin.

↳ multiple **charges**



Brings **potential cooperation**.

Costs a **membrane dissolution** on the score.

Discussion and future work

Tournaments in class

- 1 Students design valkyries.
↳ group work
- 2 We run tournaments.
- 3 Students get grades.

A classic in teaching multi-agent systems.

Teaching and research benefits

- 1 Design P systems.
- 2 Revise **probabilities**.
- 3 Think about survival in an **adversarial environment**.
↳ Potential for thinking about the **origins of Life**.
- 4 Promote various **formalisms and simulation engines**.
↳ cP systems, kernel P systems, numerical P systems, spiking neural P systems, etc.

Scoring

$$\text{score}(\Pi_j) = \frac{1}{|\Pi_j|} \left(|\Pi_j| - \frac{1}{M} \sum_{i=1}^M \text{diss}_i(\Pi_j) \right)$$

Other scoring functions?

- Better capture the **results**.
- **Avoid trivial** edge cases.
- Measure the **production** of certain symbols?
↳ forget dissolution

Core Wars vs. Queens of the Hill

Core Wars

Queens of the Hill

Data Secondary

Important

Erasure Instruction-based

Dissolution

Programs Mutable

Immutable (fixed rules)

Determinism Deterministic

Non-deterministic
(probabilistic)

$$\mathbb{Q} = \boxed{\Pi_1 \quad \dots \quad \Pi_m}$$

$$\boxed{a \rightarrow a(\delta_t, out)$$

$$a$$

$$\boxed{a \rightarrow a\bar{\delta}_t(\delta_t, out)$$

$$a$$

$$\boxed{a \rightarrow a\bar{\delta}_1^{r(m-1)}$$

$$a\bar{\delta}_1^{r(m-1)}$$

$$\boxed{a \rightarrow a(b, out)$$

$$a$$

$$\{a \rightarrow (a, out) \mid a \in \Sigma\}$$

 δ

 Subvalkyrie₁

 Subvalkyrie₂

 Subvalkyrie₃

Thank you BWMC organizers!